



SmartSpec: End-to-end specification of Smart Contract and Services

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Why verify end-to-end?

- Smart Contracts are high-value targets
 - > Large liability / small area
 - > Attack incentive / reduced cost
- Verification is a gateway to acceptance
 - > Commercial argument (Tezos, Cardano, ...)
 - > Regulatory requirement (Concordium,...)
 - Regulators may step in as soon as verification is feasible (RGPD)
- Smart Contracts are but one (key) component implementing *services*
 - > Asset management, booking, payment/delivery, directory,...
 - > Alongside non-chain code (apps, plugins), third parties (oracles), party obligations (contracts, regulations)
- Contracting parties are really concerned with end-to-end execution of services
 - > Code-only unit specifications overly technical (e.g., protocol description)

Technical challenges

- Specifications cover both Smart Contracts and general soft/hard/wet – ware
 - > Need transparent specifications for non-code components
 - > General expressiveness, unhindered by implementation constraints (with Smart Contract subset)
- Reasoning about concurrency and distribution
 - > Handle shared and fragmented state
 - > Distributed actions \rightsquigarrow synchronous execution
 - > Temporal modalities
- Reasoning about flows of information, trust, cash, gas, ...
 - > Belief modalities (BAN logic), source/authority modalities (IF logic)
- Probabilistic correctness
 - > Crypto (sub) protocols, secrecy/privacy
- Economic correctness
 - > Game theory (e.g., auctions, escrow schemes)

Technical assets

- Scilla
 - > Award-winning smart contract language design, adhering to a minimal transition model
 - > Asynchronous, implementation-oriented execution model
- TLA
 - > Turing award winning abstract specification language with synchronous transitions and time modalities
 - > Based on set theory (less suited to functional and data-oriented programming)
- Coq
 - > Fully general proof system with support for developing and embedding arbitrary theories
- MathComp
 - > Modular Coq theory library for general mathematics and combinatorics (though no crypto/game theory)
- Michelson and its compilers
 - > Execution for code fragments, simulation in general
 - > Leverage existing compilers (LIGO, SmartPy, Morley, ...)

Scilla

From ERC223 sample

(*spender: address which owner wants to give access to transfer tokens*)

transition approve (spender : ByStr20, value : Uint128)

sender_map <- allowed[_sender];

match sender_map **with**

| Some m =>

allowed[_sender][spender] := value;

e = { _eventname : "spenderapproved"; _recipient : _sender; _amount : Uint128 0; message : "spender approved" };

event e

| None =>

allowed[_sender][spender] := value;

e = { _eventname : "spenderapproved"; _recipient : _sender; _amount : Uint128 0; message : "spender approved" };

event e

end

TLA

From Dijkstra ring termination sample

$\text{SendMsg}(i) ==$
 $\wedge \text{active}[i]$
 $\wedge \exists j \in \text{Nodes} \setminus \{i\} :$
 $\wedge \text{active}' = [\text{active} \text{ EXCEPT } ![j] = \text{TRUE}]$
 $\wedge \text{color}' = [\text{color} \text{ EXCEPT } ![i] = \text{IF } j > i \text{ THEN "black" ELSE @}]$
 $\wedge \text{UNCHANGED} \langle\langle \text{tpos}, \text{tcolor} \rangle\rangle$

$\text{Deactivate}(i) ==$
 $\wedge \text{active}[i]$
 $\wedge \text{active}' = [\text{active} \text{ EXCEPT } ![i] = \text{FALSE}]$
 $\wedge \text{UNCHANGED} \langle\langle \text{color}, \text{tpos}, \text{tcolor} \rangle\rangle$

$\text{Environment} == \exists i \in \text{Nodes} : \text{SendMsg}(i) \vee \text{Deactivate}(i)$

Workplan

- Language Design
 - > Combine features of Scilla (transition monad), TLA (synchronous transition) with logical assertions/refinements making use of extended theories.
 - > Select features and semantics to balance expressiveness with formal verification feasibility.
 - > Collaboration (visiting position) with Ilya Serguei (Scilla)
- Evaluation and illustration
 - > Develop examples that test, inform, and showcase the language design decisions.
- Verification support
 - > Embed the language semantics in Coq, using MathComp; embed or compile the syntax
 - > Extend MathComp with the required relevant theories, such as concurrency with communication and synchronisation, information flow and modalities, game semantics, and probabilities and expectations.
- Implementation
 - > Compilation of smart contract code specification to Michelson (via a higher level language).
 - > Simulation of non-code specification fragments.
 - > Compilation of assertions to verification conditions for model and conformance-checking tools.

Merci !

Any questions?