

Disco: Bridging Synchrony and Asynchrony

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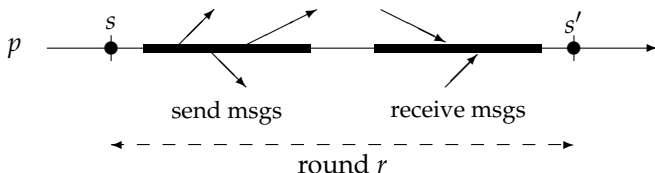
(project involving Antique and VeriDis)



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Communication-closed Rounds

- Round-based distributed algorithms



- ▶ process-local organization of computation
 - ▶ state s' computed from s and messages received
 - ▶ messages sent in round r are received only in round r
- Purely logical structure of algorithms
 - ▶ processes execute fully asynchronously
 - ▶ most classical consensus algorithms fall into this class

[Elrad & Francez 1982], [Gafni 1998], [Charron-Bost & Schiper 2006]

Reduction to Synchronous Execution Model

- Processes execute each round¹ atomically



- Rearrange rounds of different processes
 - round rnd_p^m right-commutes with rnd_q^n if $m > n$
 - messages sent during rnd_q^n do not influence rnd_p^m

¹shifted by “half a round”

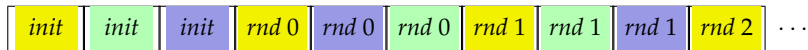
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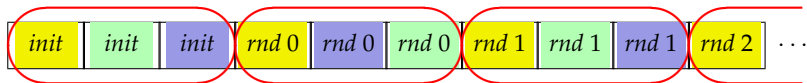
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- Executions of same round by different processes are independent
 - ▶ pretend that rounds are the unit of atomicity

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Theorem

For every asynchronous execution of a round-based distributed algorithm there exists a synchronous execution such that every process sees the same sequence of local states.

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Corollary

Any stuttering-invariant property that depends only on local views can be verified over synchronous executions.

- “Local” LTL-X formulas
 - ▶ any formula containing state variables of a single process
 - ▶ first-order combinations of “local” formulas
 - ▶ examples: distributed consensus, k -set agreement, ...

Existing Work (Antique / Veridis)

- Mechanized verification

- ▶ representation of Heard-Of model in Isabelle/HOL
- ▶ SMT-based verification of consensus algorithms
- ▶ classical algorithms tolerating crash and Byzantine failures

- PSync: domain-specific language

- ▶ express algorithms and verify them in the synchronous model
- ▶ generate executables for asynchronous execution
- ▶ strong guarantees and reasonable performance

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 - ▶ support full Byzantine failures, including signed messages
 - ▶ take into account probabilistic guarantees
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- 3 Design an execution platform
 - ▶ efficient execution of programs in an asynchronous context
 - ▶ account for timeouts, digital signatures, optimizations
 - ▶ formally prove correctness in a proof assistant