

Plan

- Why3
- demos
- conclusions

Goal

Write elegant programs

with elegant correctness proofs

+ training in program proofs



Why3 (1/8)

A programming language tells you **what** a program does, Why3 tells you **why** it works.

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- 3rd release of system Why
- developed at LRI (orsay) + Inria
- http://why3.lri.fr

[Jean-Christophe Filliâtre, Claude Marché, Andrei Paskevich, Guillaume Melquiond, Vincent Bolot, et al]

Why3 (2/8)

small Pascal-like imperative programming language

[with ML syntax 🙁 !!]

• invariants + assertions in Hoare logic

[+ recursive functions, inductive datatypes, inductive predicates]

• interfaces with modern SMT's

[alt-ergo, cvc3, cvc4, eprover, gappa, simplify, spass, yices, z3]

• interfaces with interactive proof assistants

[coq, pvs, isabelle-hol]

Why3 (3/8)

```
    programming language MLW

    let swap (a: array int) (i: int) (j: int) =
    let v = a[i] in
      a[i] <- a[j];
      a[j] <- v
    let selection_sort (a: array int) =
      for i = 0 to length a - 1 do
        let imin = ref i in
        for i = i + 1 to length a - 1 do
          if a[j] < a[!imin] then imin := j</pre>
        done;
        swap a !imin i
      done
                           imin
   0
                   i
а
```

Why3 (4/8)

```
    Hoare logic
```

```
let swap (a: array int) (i: int) (j: int) =
 let v = a[i] in
  a[i] <- a[j];</pre>
  a[i] <- v
let selection_sort (a: array int) =
   for i = 0 to length a - 1 do
    let imin = ref i in
    for j = i + 1 to length a - 1 do
      invariant { i <= !imin < i }</pre>
      invariant { forall k: int. i <= k < j \rightarrow a[!imin] <= a[k] }
      if a[j] < a[!imin] then imin := j
    done;
    swap a !min i
  done
   0
                               imin
а
```

Why3 (5/8)

theories on arrays

```
let swap (a: array int) (i: int) (j: int) =
  requires { 0 <= i < length a / 0 <= j < length a }
  ensures { exchange (old a) a i j }
  let v = a[i] in
  a[i] <- a[j];
  a[j] <- v</pre>
```

(see the why3 libraries)

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http://why3.lri.fr

Why3 (6/8)

theories on arrays let selection_sort (a: array int) = ensures { sorted a \land permut (old a) a } 'L: for i = 0 to length a - 1 do invariant { sorted_sub a 0 i \land permut (at a 'L) a} invariant { forall k1 k2: int. $0 \le k1 \le i \le k2 \le length a \rightarrow a[k1] \le a[k2]$ } let imin = ref i in for j = i + 1 to length a - 1 do invariant { i <= !imin < j }</pre> invariant { forall k: int. i <= k < j -> a[!imin] <= a[k] } if a[j] < a[!imin] then imin := j</pre> done: swap a !imin i ; done imin а

Why3 (8/8)

- interfaces with interactive proof assistants
- PVS [SRI, Shankar], Isabelle [Paulson, Nipkow]
- Coq [Inria, Herbelin et al]
 - Why3 theories are translated to Coq
 - lengthy proofs are feasible
 - use Ssreflect commands to shorten proofs [MSR-Inria, Gonthier
 et al]
 - unfortunately Why3 is not fully compatible with SSreflect

Why3 (7/8)

- interfaces with automatic provers (SMT's)
- SMT tool successful if «good assertion»
- impact on writings of Hoare logic formulae
- impact on program text
- Alt-Ergo among best for Why3 [LRI, Conchon, et al]
- Z3 is excellent [MSRR, Bjorner/de Moura]
- CVC3 top on recursive datatypes
- Gappa for real numbers [Inria, Melquiond]

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A few sorting algorithms

• demos

• insertion sort



Depth-first search in graphs

- reachability [the 'white path theorem']
- non white-to-black edges in undirected graphs



- acyclicity test
- articulation point
- strongly connected components

A few sorting algorithms

quicksort







Conclusion (1/3)

- Automatic part of proof for tedious case analyzes
- Interactive proofs for the conceptual part of the algorithm
- the ideal world
- From interactive part, one must call the automatic part
 - possible extensions of Why3 theories
 - but typing problems (inside Coq)

Conclusion (3/3)

- Why3 is excellent for mixing formal proofs and SMT's calls
- Interface still rough for beginners
- Concurrency ?
- Functional programs ?
- Hoare logic vs Type refinements (F* [MSR])
- Frama-C project at french CEA extends Why3 to C programs.

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Conclusion (2/3)

- Hoare logic prevents to write awkward denotational semantics
- Nobody cares about termination ?!



- Explore simple programs about algorithms before jumping to large programs.
- Why3 memory model is naive. It is a «back-end for other systems».
- Plan to experiment on graph algorithms and prove all Sedgewick's book on algorithms.

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